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# A Security Policy Definition Language, Semantics, and Open Source Tools

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# Overview

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How can we concisely and uniformly express fine-grained security policies?

The plan in three parts:

1. Introduction to our domain-specific language (DSL) for security policies.
2. DSL simulation tools for RISC-V.
3. DSL hardware enforcement on RISC-V.

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# 1) A Security Policy DSL

# Example: A policy for RWX permissions

```
metadata:  
  Rd | Wr | Ex
```

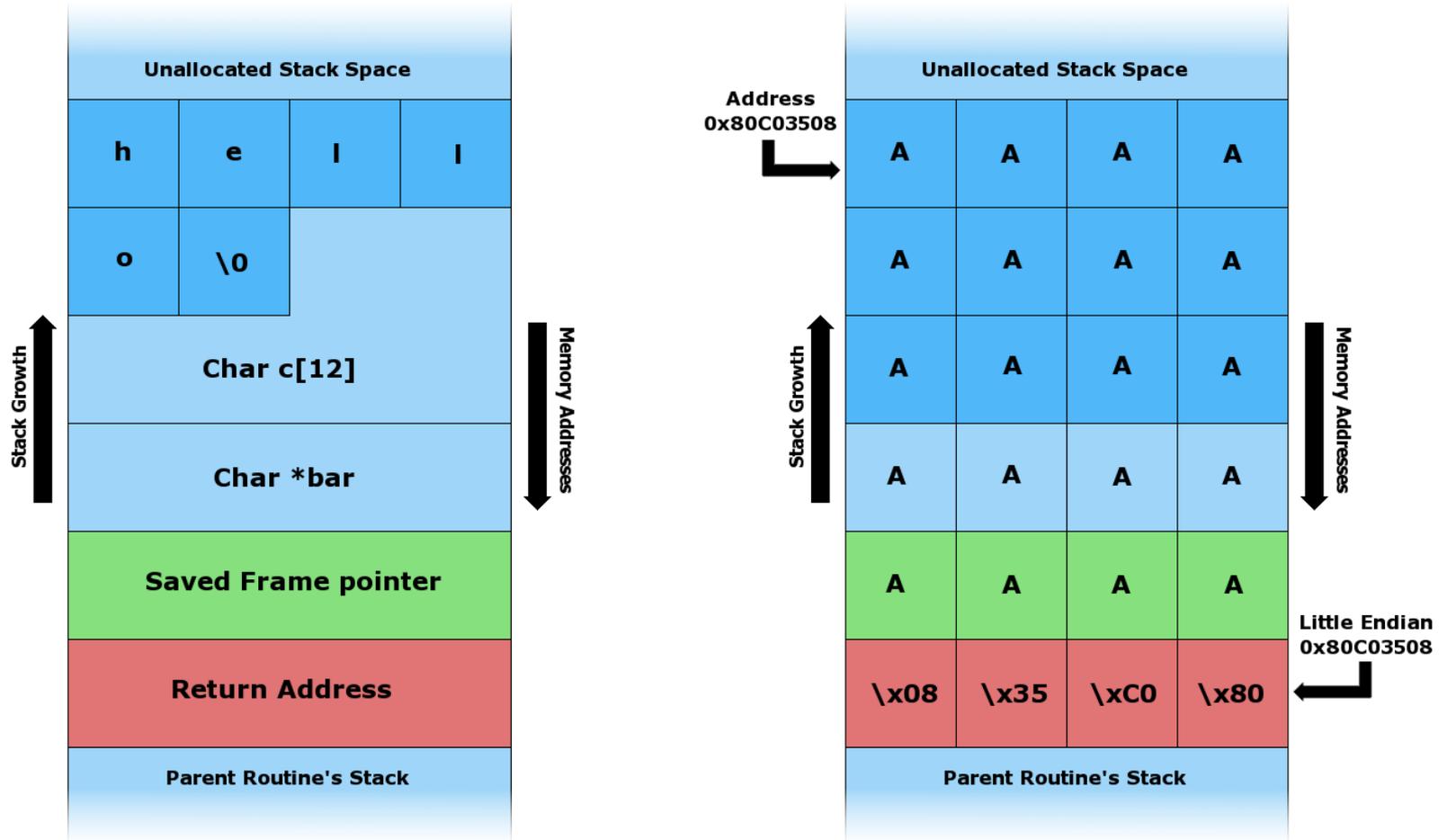
```
policy:
```

```
  rwxPol =  
    loadGrp ( code == [+Ex], env == _, mem == [+Rd]  
             -> env = env)  
  ^ storeGrp ( code == [+Ex], env == _, mem == [+Wr]  
             -> mem = mem, env = env)  
  ^ nonMemGrp ( code == [+Ex], env == _  
               -> env = env)
```

A policy is a list of rules. Instructions not matched are rejected.

```
require:  
  init Link.MemoryMap.UserStack = {Rd, Wr}  
  init Link.MemoryMap.UserHeap  = {Rd, Wr}  
  init Elf.Section.Code         = {Ex}
```

# Example: Stack Buffer Overflow



# Stack Protection Policy

- Tag memory containing function prologue and epilogue code.
- No other code may touch stack frame.

```
metadata:
  Frame | Prologue | Epilogue

policy:
  stackPol =
    // Prologue creates stack Frame
    storeGrp (  code == [+Prologue], mem == __, env == _
              -> mem = {Frame})

    // Other code can't touch Frame
    ^ storeGrp (  code == [-Prologue, -Epilogue], mem == [+Frame]
              -> fail "Illegal stack access")

    // Epilogue clears Frame, allow other writes, ...
    ...
```

# Example: Simple CFI

- Statically determine legal jump destinations.
- Dynamically check each jump.

```
metadata:
  Target | Jumping

policy:
  stackPol =
    // jump instructions set "Jumping"
    jumpGrp (env == _ -> env = env[+Jumping])

    // Landing on legal target clears "Jumping"
    ^ allGrp ( code == [+Target], env == [+Jumping]
              -> env = env[-Jumping])

    // Landing elsewhere is a violation
    ^ allGrp ( code == [-Target], env == [+Jumping]
              -> fail "Illegal jump")

    ^ ...
```

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## **2) DSL Simulation Tools for RISC-V**

# Open Source Simulation Tools

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- Our simulator runs RISC-V programs with policy enforcement.
- Built on QEMU, plugin does policy checks.
- Many policies and test programs.
  - FreeRTOS webserver with application-specific policies.
  - Initial seL4 support.
- Open-source, MIT license, we'd love collaborators:
  - <https://github.com/draperlaboratory/hope-tools>

# RISC-V Emulation with Policies

- Simulator detects policy violations in application.
- Provides user feedback on location and cause of violation.

```
cjc1527@leopard> cat pex.log
Policy Violation:
  PC = 20451ad0      MEM = 80502dd8
Metadata:
  Env   : {}
  Code  : {storeGrp}
  Op1   : {}
  Op2   : {}
  Op3   : -0-
  Mem   : {Frame}
Explicit Failure
Illegal stack access
MSG: End test.
cjc1527@leopard> █
```

# Metadata Support in GDB

- Simulator supports GDB.
- New commands to inspect metadata as application runs.

```
(gdb) metadata
Metadata related commands:
  pvm          - print violation message
  env-m        - get the env metadata
  reg-m n      - get register n metadata
  csr-m a      - get csr metadata at addr a
  mem-m a      - get mem metadata at addr a
Watchpoints halt simulation when metadata changes
  env-mw       - set watch on the env metadata
  reg-mw n     - set watch on register n metadata
  csr-mw a     - set watch on csr metadata at addr a
  mem-mw a     - set watch on mem metadata at addr a
(gdb) mem-m 0x20417688
{loadGrp, notMemGrp}
(gdb) █
```

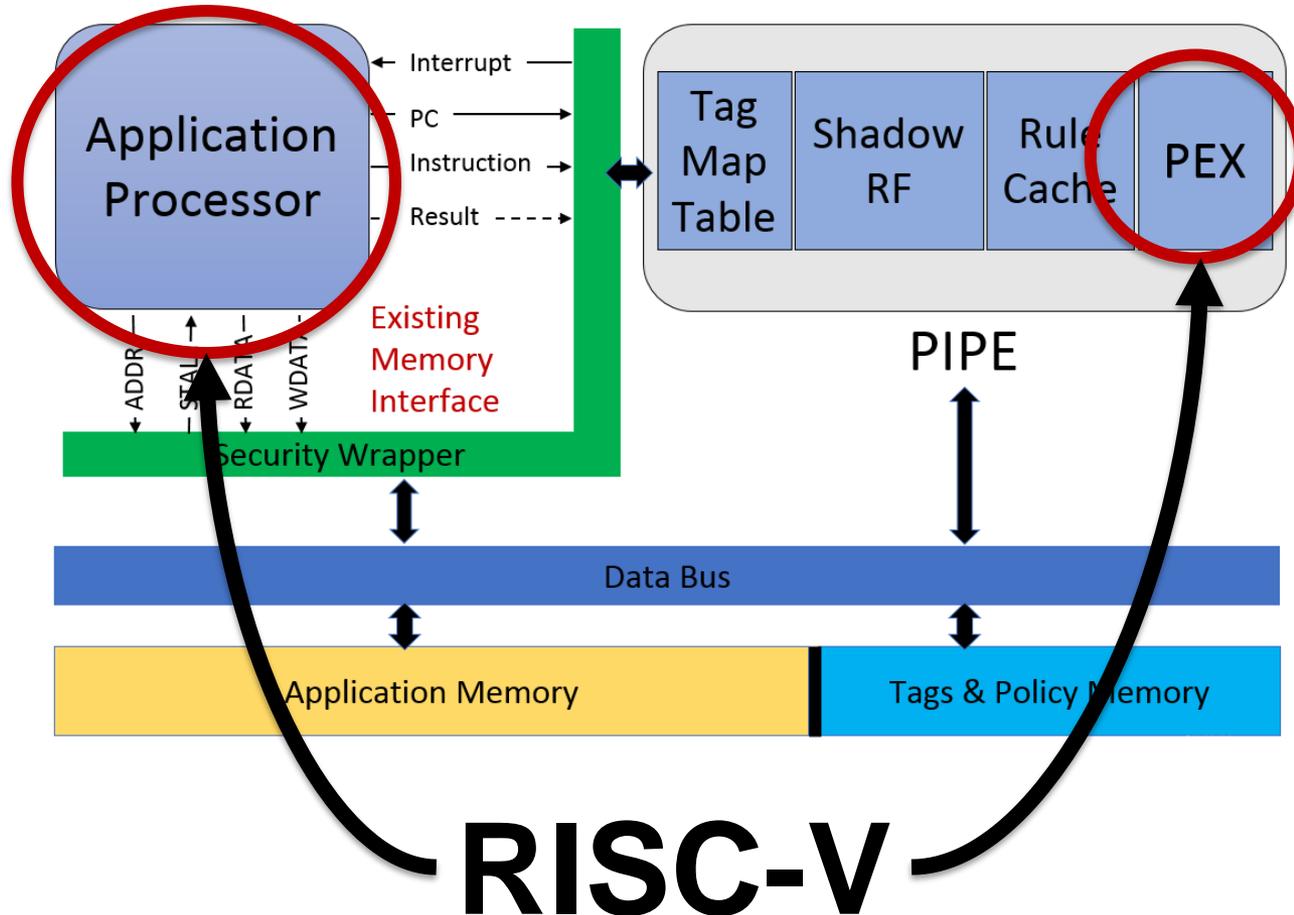
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## **3) DSL Hardware Enforcement on RISC-V**

# PIPE Hardware Diagram

- PIPE: Processor Interlocks for Policy Enforcement



# HW Availability

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- Dover and Draper have collaborated on the design of PIPE.
- Dover offers a performant, commercial version, called CoreGuard.



**COREGUARD<sup>®</sup>**  
Protected. Trusted.

- Draper integrates PIPE for government customers in defense systems.

# Conclusion

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- Summary:
  - DSLs can make security policies easier to write, understand, and verify.
  - You can play with our open simulation tools today:
    - <https://github.com/draperlaboratory/hope-tools>
  - Talk to us about performant HW implementations for RISC-V (and other architectures)
  - Greg Sullivan : [gregs@dovermicrosystems.com](mailto:gregs@dovermicrosystems.com)
  - Chris Casinghino: [ccasinghino@draper.com](mailto:ccasinghino@draper.com)

Thanks!

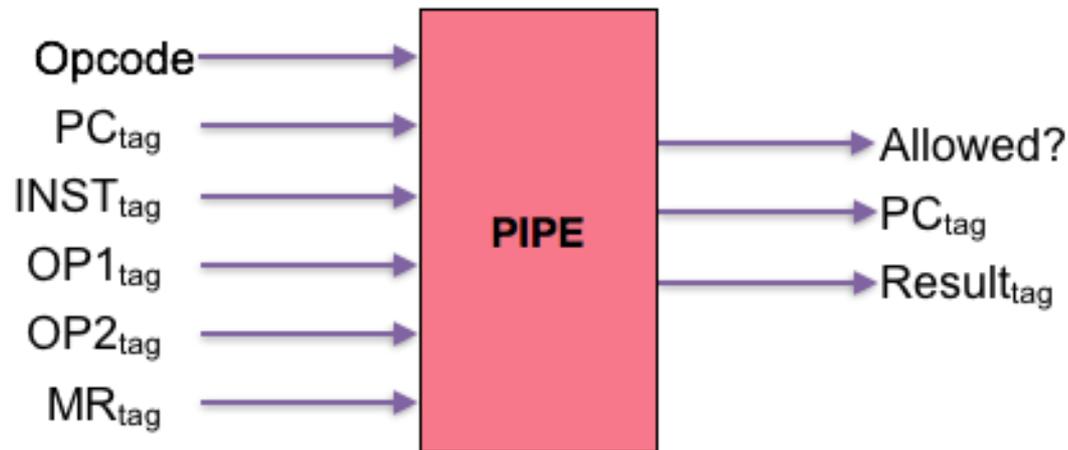
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**Backup**

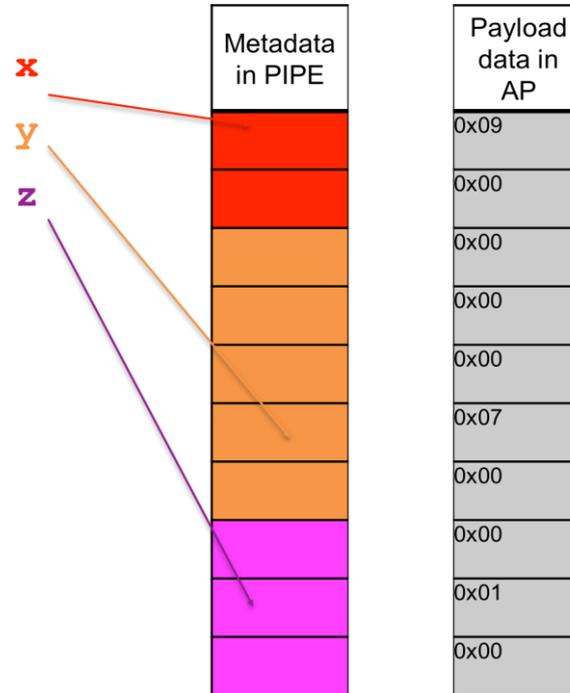
# PIPE hardware

- PIPE: Processor Interlocks for Policy Enforcement
- Runs in parallel with application processor.
- Generates application processor interrupt if policy violation occurs.
- Creates unassailable hardware interlock blocking execution of bad instructions.
- Cache keeps performance high.



# Example: Heap Memory Safety

- **Policy Goal:** enforce spatial and temporal safety
- **Method:** “color” pointers and memory
  - On allocation, add colors to the region and the pointer.
  - Recolor on free.
- **Policy:**
  - On access, pointer and memory colors must match.



```
z = malloc(3);  
z[1] = 0x01;  
y = malloc(5);  
y[3] = 0x07;  
x = malloc(2);  
x[0] = 0x09;  
x[2] = 0xbad; //FAIL
```

# Heap Safety Policy

- Tags can carry data values, allowing us to implement this color scheme.

```
type:
  data Color = Int (20)

metadata:
  ApplyColor | NewColor | Pointer Color | Cell Color | ...

policy:
  // Generate new color during malloc
  storeGrp(   code == [ApplyColor], mem == [NewColor]
             -> mem = mem[(Pointer new)])

  // Allow load if colors match
  ^ loadGrp(  mem == [Cell color], addr == [Pointer color]
            -> res = mem[-(Cell _)])

  ...
```

# Policy Language Continued

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- Users may define new “opgroups”.
  - Group instructions in the way that makes sense for your policy.
  - Keeps policy rules somewhat ISA-agnostic.
  
- Language implemented in Haskell.
  - We generate C and compile to RISC-V binaries.
  - Lots of room left to optimize.
  
- Language has formal operational semantics.
  - Future goal: prove implementation correct.